

and  $\deg(e_2(x)) \geq m_0$ . If  $f \leq m_0$ , then  $x^i e_1(x) \not\equiv e_2(x) \pmod{g_1(x^b)g_2(x)}$ . Therefore the shortened cyclic code generated by  $g(x) = g_1(x^b)g_2(x)$  is capable of correcting all single-burst errors of length  $\min\{m_0, bt\}$  or less.

Earlier we presented a class of shortened cyclic codes for a compound channel [4]. A code in this class was generated by the polynomial  $g(x) = (x^d + 1)g'(x)$  where  $g'(x)$  generates a  $t$ -error-correcting cyclic code of length  $N$ . A comparison shows that for the same number of check digits, the shortened cyclic codes presented here have a better rate of transmission, but the earlier shortened cyclic codes have a better guaranteed burst-error correction capability.

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## On the Fractional Weight of Distinct Binary $n$ -Tuples

JAMES L. MASSEY

**Abstract**—It is shown that the fraction  $p$  of ones in the  $Mn$  positions of  $M$  distinct binary  $n$ -tuples satisfies the inequality

$$h(p) \geq (1/n) \log_2 M$$

where  $h(p) = -p \log_2 p - (1-p) \log_2 (1-p)$  is the binary entropy function. This inequality, which simplifies the derivation of the distance property of the Justesen codes, is proved using an elegant information-theoretic argument due to Kriz.

Let  $p$  be the fraction of the  $Mn$  positions in  $M$  distinct binary  $n$ -tuples which contain ones, and let  $p_i$  be the fraction of these  $n$ -tuples whose  $i$ th component is a one; then

$$p = \frac{p_1 + p_2 + \cdots + p_n}{n}$$

We now derive an inequality relating  $p$  and  $n$  using an information-theoretic argument employed by Kriz [1] for a closely related problem.

Let  $[x_1, x_2, \cdots, x_n]$  be a random  $n$ -tuple, which takes on each of the  $M$  given distinct  $n$ -tuple values with probability  $1/M$ . Then the uncertainty of this random  $n$ -tuple is

$$H(X_1 X_2 \cdots X_n) = \log_2 M.$$

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The author is with the Department of Electrical Engineering, University of Notre Dame, Notre Dame, Ind. 46556.

As always the joint uncertainty of  $n$  letters is overbounded by the sum of their individual uncertainties, i.e.,

$$H(X_1 X_2 \cdots X_n) \leq H(X_1) + H(X_2) + \cdots + H(X_n)$$

and because of our probabilistic assignment

$$H(X_i) = h(p_i), \quad 1 \leq i \leq n$$

where  $h(p) = -p \log_2 p - (1-p) \log_2 (1-p)$ . Combining these expressions, we have

$$h(p_1) + h(p_2) + \cdots + h(p_n) \geq \log_2 M$$

and also

$$\begin{aligned} h(p_1) + \cdots + h(p_n) &= n \left[ \frac{1}{n} h(p_1) + \cdots + \frac{1}{n} h(p_n) \right] \\ &\leq nh \left( \frac{p_1 + p_2 + \cdots + p_n}{n} \right) = nh(p) \end{aligned}$$

where we have exploited the convexity of  $h$ . We have thus proved the following theorem.

**Theorem:** The fraction  $p$  of ones in the  $Mn$  positions of  $M$  distinct binary  $n$ -tuples satisfies

$$h(p) \geq \frac{1}{n} \log_2 M.$$

It is interesting to note that the bound of this theorem is tight for a)  $M = 1$ , since  $h(p) \geq 0$  implies  $0 \leq p \leq 1$  and the values  $p = 0$  and  $p = 1$  are achievable, and for b)  $M = 2^n$ , since  $h(p) \geq 1$  implies the unique solution  $p = \frac{1}{2}$ . Moreover, the bound is surprisingly tight in general. For instance, with  $n = 20$  and  $M = 2^{10} = 1024$ , the bound gives  $h(p) \geq \frac{1}{2}$ , which implies  $0.110 \leq p \leq 0.890$ . The actual minimum and maximum achievable values of  $p$  for this case are 0.139 and 0.861, respectively.

The reader will recognize the theorem as a strengthening and simplification of the lemma of Justesen [2], which was the key inequality used in the derivation of the asymptotic distance bound for the Justesen codes.

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## On Binary Majority-Logic Decodable Codes

S. G. S. SHIVA AND S. E. TAVARES

**Abstract**—Let  $V'$  be a binary  $(n, k)$  majority-logic decodable code with  $g'(X)$  as its generator polynomial and odd minimum distance  $d$ . Let  $V$  be the  $(n, k-1)$  subset code generated by  $g'(X)(1+X)$ . This correspondence shows that  $V$  is majority-logic decodable with  $d+1$  orthogonal estimates. This fact is useful in the simultaneous correction of random errors and erasures.

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S. G. S. Shiva is with the Electrical Engineering Department, University of Ottawa, Ottawa, Ont., Canada.

S. E. Tavares is with the Electrical Engineering Department, Queen's University, Kingston, Ont., Canada.