



## Exercise 4 of March 16, 2017

<http://www.isi.ee.ethz.ch/teaching/courses/it2.html>

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### Problem 1

### *“Double”-Gaussian Channel*

- Consider a discrete-time memoryless Gaussian channel with average-power constraint  $E_s = 1$ . Sketch the capacity of this channel as a function of the noise variance  $N$ .
- You are given two parallel independent Gaussian channels with noise variances  $N_1$  and  $N_2$ , respectively. You are allocated a total power of  $E_s$  to use between these two channels. The sum of the noise variances in the two channels is  $2N$ , but you will be able to choose from the two schemes below as to how the noise is distributed ( $\Delta$  and  $N$  are fixed).

**Scheme A:**  $N_1 = N_2 = N$ ,

**Scheme B:**  $N_1 = N + \Delta$ ,  $N_2 = N - \Delta$ ,  $0 < \Delta < N$ .

- If you must put half of your power allocation in each channel, which scheme would you prefer?
- If you are allowed to arbitrarily divide your power allocation between the two channels, which scheme would you prefer?

### Problem 2

### *Bandlimited Gaussian Channel*

Consider the bandlimited Gaussian channel with noise power spectral density  $N_0/2$  and power  $P$ . The capacity of this channel is

$$W \log \left( 1 + \frac{P}{N_0 W} \right) \text{ bits per second,}$$

where  $W$  is the bandwidth of the channel. What would you rather have, twice the bandwidth or twice the power?

### Problem 3

### *Exponential Noise Channel*

Consider the channel

$$Y_k = X_k + Z_k,$$

where the inputs  $\{X_k\}$  are nonnegative and where  $\{Z_k\}$  are IID exponentially distributed with mean  $\mu$ . Assume that the inputs satisfy the mean constraint

$$E[X_k] \leq \lambda, \quad k = 0, 1, \dots$$

Show that the capacity of this channel is

$$C = \log \left( 1 + \frac{\lambda}{\mu} \right).$$