



Exercise 7 of April 6, 2017

<http://www.isi.ee.ethz.ch/teaching/courses/it2.html>

Problem 1

Cooperation on the Gaussian MAC

Derive the capacity region of the Gaussian MAC with cooperation. More specifically, derive the capacity region of the MAC

$$Y = X_1 + X_2 + Z,$$

where both transmitters have access to both messages, $Z \sim \mathcal{N}(0, \sigma^2)$, and for every pair of messages, the average power of X_1 must not exceed E_1 and the average power of X_2 must not exceed E_2 .

Problem 2

Gaussian MAC

Consider a Gaussian MAC

$$Y = X_1 + X_2 + Z,$$

where $X_1 \in \mathbb{R}$ is the input from transmitter 1, $X_2 \in \mathbb{R}$ is the input from transmitter 2, both transmitters are independent, Z is a Gaussian random variable with mean zero and variance σ^2 , and Y is the channel output. We impose average-power constraints on both inputs, i.e., the average power of every codeword from transmitter 1 must not exceed E_1 and the average power of every codeword from transmitter 2 must not exceed E_2 . Show that the capacity region of this channel is the set of all rate pairs (R_1, R_2) satisfying

$$\begin{aligned} R_1 &\leq \frac{1}{2} \log \left(1 + \frac{E_1}{\sigma^2} \right), \\ R_2 &\leq \frac{1}{2} \log \left(1 + \frac{E_2}{\sigma^2} \right), \\ R_1 + R_2 &\leq \frac{1}{2} \log \left(1 + \frac{E_1 + E_2}{\sigma^2} \right). \end{aligned}$$

Problem 3

Gaussian MAC and TDMA

Suppose that on a Gaussian MAC with equal average-power constraints $E_1 = E_2 = E$ a rate pair (R, R) is in the interior of the capacity region. Show that this rate pair can be achieved by *time-division multiplexing*, where at each channel use, at most one input of the MAC is nonzero.