



Exercise 11 of May 11, 2017

<http://www.isi.ee.ethz.ch/teaching/courses/it2.html>

Problem 1

Method of Types

Let $\mathcal{X} = \{a, b, c\}$ be a ternary alphabet. Consider sequences $\mathbf{x} \in \mathcal{X}^5$ of length $n = 5$.

- What is the exact number of types that \mathbf{x} can take on, i.e., what is $|\mathcal{P}_5|$? Compare this number to the upper bounds derived in class.
- What is the type of $\mathbf{x} = (b, c, c, a, b)$?
- How many sequences have the same type as $\mathbf{x} = (b, c, c, a, b)$? Compare this number to the upper and the lower bounds derived in class. Are these bounds useful? Is there an easy way to improve the lower bound?
- Assume that \mathbf{X} is drawn IID Q where $Q(a) = 2/3$ and $Q(b) = Q(c) = 1/6$. What is the probability that $\mathbf{X} = (b, c, c, a, b)$? Compute it both directly and using types.
- Still assuming that \mathbf{X} is drawn IID Q , give an upper and a lower bound to the probability that $P_{\mathbf{X}} = P_{(b,c,c,a,b)}$. Can you compute this probability exactly?

Problem 2

Error Exponent for Universal Codes

A universal code of rate R applied to an IID source of distribution Q achieves an error probability

$$P_e^{(n)} \leq (n+1)^{|\mathcal{X}|} e^{-nD(P^*||Q)},$$

where P^* achieves $\min D(P||Q)$ over all P such that $H(P) \geq R_n$.

- For given Q and R_n , what form does P^* have?
Hint: You may use Lagrange multipliers to solve this part.
- Let X be binary. Find the region of source distributions Q for which rate R is sufficient for the universal source code to achieve $P_e^{(n)} \rightarrow 0$.

Problem 3

Large Deviations

Let X_1, X_2, \dots, X_n be IID $\mathcal{N}(0, \sigma^2)$. Find the exponent in the probability that

$$\frac{1}{n} \sum_{i=1}^n X_i^2 \geq \alpha^2.$$

Hint: Although the alphabet is infinite, you may use Sanov's theorem to solve this problem.